

(MS-)CHAPv2

Challenge/Response Authentication Protocol, version 2

Protocol Purpose

Mutual authentication between a server and a client who share a password. CHAPv2 is the authentication protocol for the Point-to-Point Tunneling Protocol suite (PPTP).

Definition Reference

[Zor00]

Model Authors

- Haykal Tej, Siemens CT IC 3, 2003
- Paul Hankes Drielsma, ETH Zürich

Alice&Bob style

We assume that the server B and client A share password $k(A, B)$ in advance. The server and client generate nonces Nb and Na , respectively.

1. $A \rightarrow B : A$
2. $B \rightarrow A : Nb$
3. $A \rightarrow B : Na, H(k(A, B), (Na, Nb, A))$
4. $B \rightarrow A : H(k(A, B), Na)$

Model Limitations

Issues abstracted from:

- Message structure: As is standard, we abstract away from the concrete details of message structure such as bit lengths, etc. What is left after this abstraction contains several redundancies, however (at least in the Dolev-Yao model). We therefore eliminate these redundancies, retaining the core of the data dependencies of the protocol.

Problems considered: 3

Attacks Found

None

Further Notes

A cryptanalysis of this protocol in its full complexity can be found in [SMW99].

HLP SL Specification

```
role chap_Init (A,B      : agent,
                 Kab     : symmetric_key,
                 H       : function,
                 Snd, Rcv: channel(dy))

played_by A
def=

local State  : nat,
      Na, Nb : text

const sec_kab1 : protocol_id

init  State := 0

transition
1. State    = 0 /\ Rcv(start) =|>
   State' := 1 /\ Snd(A)

2. State    = 1 /\ Rcv(Nb') =|>
   State' := 2 /\ Na' := new() /\ Snd(Na'.H(Kab.Na'.Nb'.A))
   /\ witness(A,B,na,Na')
   /\ secret(Kab,sec_kab1,{A,B})
```

```

3. State    = 2 /\ Rcv(H(Kab.Na)) =|>
   State' := 3 /\ request(A,B,nb,Nb)

end role

```

```

role chap_Resp (B,A : agent,
                 Kab : symmetric_key,
                 H: function,
                 Snd, Rcv: channel(dy))

played_by B
def=

local State  : nat,
      Na, Nb : text

const sec_kab2 : protocol_id

init  State := 0

transition
1. State    = 0 /\ Rcv(A') =|>
   State' := 1 /\ Nb' := new() /\ Snd(Nb')
   /\ witness(B,A,nb,Nb')

2. State    = 1 /\ Rcv(Na'.H(Kab.Na'.Nb.A)) =|>
   State' := 2 /\ Snd(H(Kab.Na'))
   /\ request(B,A,na,Na')
   /\ secret(Kab,sec_kab2,{A,B})

end role

```

```

role session(A,B: agent,
            Kab: symmetric_key,
            H: function)
def=

local SA, SB, RA, RB: channel (dy)

composition

```

```
chap_Init(A, B, Kab, H, SA, RA)
/\ chap_Resp(B, A, Kab, H, SB, RB)
end role
```

```
role environment()
def=

const a, b          : agent,
kab, kai, kbi : symmetric_key,
h          : function,
na, nb       : protocol_id

intruder_knowledge = {a, b, h, kai, kbi }

composition
    session(a,b,kab,h) /\ 
    session(a,i,kai,h) /\ 
    session(b,i,kbi,h)

end role
```

```
goal

%secrecy of the shared key
secrecy_of sec_kab1, sec_kab2

%CHAP_Init authenticates CHAP_Resp on nb
authentication_on nb
%CHAP_Resp authenticates CHAP_Init on na
authentication_on na

end goal

environment()
```

References

- [SMW99] Bruce Schneier, Mudge, and David Wagner. Cryptanalysis of microsoft's PPTP authentication extensions (MS-CHAPv2). In *CQRE: International Exhibition and Congress on Secure Networking – CQRE [Secure]*, 1999.
- [Zor00] G. Zorn. RFC 2759: Microsoft PPP CHAP Extensions, Version 2, January 2000. Status: Informational.