# TD1 MVFA: modeling concurrent systems

## Exercise 1

Let  $TS = (S, Act, \rightarrow, I, AP, L)$  be a transition system.

- TS is called *action-deterministic* if  $|I| \leq 1$  and  $|Post(s, \alpha)| \leq 1$  for all states s and actions  $\alpha$ .
- TS is called AP-deterministic if  $|I| \leq 1$  and  $|Post(s) \cap \{s' \in S \mid L(s') = A\}| \leq 1$  for all states  $s \in S$  and  $A \in 2^{AP}$ .

Consider the following transition system  $TS_1$ .



- 1. Give the formal definition of  $TS_1$ .
- 2. Specify a finite and an infinite execution of  $TS_1$ .
- 3. Show whether  $TS_1$  is an AP-deterministic or an action-deterministic transition system.

## Exercise 2

We consider the handshaking operator  $||_H$  with H a set of actions. The definition is the same than the interleaving operator, except for the transition function for actions  $\alpha \in H$  in  $TS_1||_H TS_2$ : instead of just taking one of the transitions in  $TS_1$  or  $TS_2$  by  $\alpha$ , both transition systems must take a transition by  $\alpha$  simultaneously. For actions  $\alpha \notin H$ , the definition of the transition function is the same than in the interleaving operator.

- 1. Formally define  $TS_1||_H TS_2$  for two transition systems  $TS_1$  and  $TS_2$ .
- 2. Show that  $||_H$  is associative, that is, for any transition systems  $TS_1, TS_2, TS_3$ :

 $(TS_1||_H TS_2)||_H TS_3 = TS_1||_H (TS_2||_H TS_3)$ 

### Exercise 3

We are given two processes  $P_1, P_2$  with shared integer variable x. The program of process  $P_i$  is as follows:

```
for k_i = 1, \dots, N_i do

\begin{vmatrix} \text{LOAD}(x); \\ \text{INC}(x); \\ \text{STORE}(x); \end{vmatrix}

end
```

#### Algorithm 1: Process $P_i$

with  $N_i \ge 1$ . That is,  $P_i$  executes  $N_i$  times the assignment x := x + 1. The assignment x := x + 1 is realized using the three actions LOAD(x), INC(x) and STORE(x). Consider the parallel program:

 $\begin{aligned} x &:= 0\\ P_1 || P_2 \end{aligned}$ 

#### Algorithm 2: Parallel program P

- 1. Sketch the program graph for P with  $N_1 = N_2 = 2$ . Does P have an execution that halts with the terminal value x = 2?
- 2. What is the size of the program graph for  $N_1 = N_2 = 100$ ? For  $N_1 = N_2 = 10000$ ? Same question for the associated transition system.

#### Exercise 4

We consider the following two sequential hardware circuits:



For each  $i \in \{1, 2\}$ ,  $x_i$  is the input of Circuit  $i, y_i$  the output and  $r_i$  is a register.

The values of  $r_i$  and  $y_i$  depend on the current value of  $x_i$  and the previous value of  $r_i$ . For instance, the new value of  $r_2$  is determined by the function  $\delta_{r_2}(r_2, x_2) = r_2 \vee x_2$ , and The value of the output  $y_i$ depends on the values of  $x_i$  and  $r_i$ , for instance the value of  $y_2$  corresponds to the value of the function  $f_{y_2}(x_2, r_2) = x_2 \wedge r_2$ .

- 1. Write the functions  $f_{y_1}$  and  $\delta_{r_1}$ .
- 2. Represent Circuit 1 and Circuit 2 as transition systems.
- 3. Determine the reachable part of the transition systems of the synchronous product of these transition systems. Assume that the initial values of the registers are  $r_1 = 0$  and  $r_2 = 1$ .