## TD1 MVFA: modeling concurrent systems

## Exercise 1

Let $T S=(S, A c t, \rightarrow, I, A P, L)$ be a transition system.

- TS is called action-deterministic if $|I| \leq 1$ and $|\operatorname{Post}(s, \alpha)| \leq 1$ for all states $s$ and actions $\alpha$.
- $T S$ is called $A P$-deterministic if $|I| \leq 1$ and $\left|\operatorname{Post}(s) \cap\left\{s^{\prime} \in S \mid L\left(s^{\prime}\right)=A\right\}\right| \leq 1$ for all states $s \in S$ and $A \in 2^{A P}$.

Consider the following transition system $T S_{1}$.


1. Give the formal definition of $T S_{1}$.
2. Specify a finite and an infinite execution of $T S_{1}$.
3. Show whether $T S_{1}$ is an $A P$-deterministic or an action-deterministic transition system.

## Exercise 2

We consider the handshaking operator $\|_{H}$ with $H$ a set of actions. The definition is the same than the interleaving operator, except for the transition function for actions $\alpha \in H$ in $T S_{1} \|_{H} T S_{2}$ : instead of just taking one of the transitions in $T S_{1}$ or $T S_{2}$ by $\alpha$, both transition systems must take a transition by $\alpha$ simultaneously. For actions $\alpha \notin H$, the definition of the transition function is the same than in the interleaving operator.

1. Formally define $T S_{1} \|_{H} T S_{2}$ for two transition systems $T S_{1}$ and $T S_{2}$.
2. Show that $\|_{H}$ is associative, that is, for any transition systems $T S_{1}, T S_{2}, T S_{3}$ :

$$
\left(T S_{1} \|_{H} T S_{2}\right)\left\|_{H} T S_{3}=T S_{1}\right\|_{H}\left(T S_{2} \|_{H} T S_{3}\right)
$$

## Exercise 3

We are given two processes $P_{1}, P_{2}$ with shared integer variable $x$. The program of process $P_{i}$ is as follows:

```
for \(k_{i}=1, \ldots, N_{i}\) do
    LOAD \((x)\);
    \(\operatorname{INC}(x)\);
    STORE \((x)\);
end
```

Algorithm 1: Process $P_{i}$
with $N_{i} \geq 1$. That is, $P_{i}$ executes $N_{i}$ times the assignment $x:=x+1$. The assignment $x:=x+1$ is realized using the three actions $\operatorname{LOAD}(x), \operatorname{INC}(x)$ and $\operatorname{STORE}(x)$. Consider the parallel program:
$x:=0$
$P_{1} \| P_{2}$
Algorithm 2: Parallel program $P$

1. Sketch the program graph for $P$ with $N_{1}=N_{2}=2$. Does $P$ have an execution that halts with the terminal value $x=2$ ?
2. What is the size of the program graph for $N_{1}=N_{2}=100$ ? For $N_{1}=N_{2}=10000$ ? Same question for the associated transition system.

## Exercise 4

We consider the following two sequential hardware circuits:


Circuit 1


Circuit 2

For each $i \in\{1,2\}, x_{i}$ is the input of Circuit $i, y_{i}$ the output and $r_{i}$ is a register.
The values of $r_{i}$ and $y_{i}$ depend on the current value of $x_{i}$ and the previous value of $r_{i}$. For instance, the new value of $r_{2}$ is determined by the function $\delta_{r_{2}}\left(r_{2}, x_{2}\right)=r_{2} \vee x_{2}$, and The value of the output $y_{i}$ depends on the values of $x_{i}$ and $r_{i}$, for instance the value of $y_{2}$ corresponds to the value of the function $f_{y_{2}}\left(x_{2}, r_{2}\right)=x_{2} \wedge r_{2}$.

1. Write the functions $f_{y_{1}}$ and $\delta_{r_{1}}$.
2. Represent Circuit 1 and Circuit 2 as transition systems.
3. Determine the reachable part of the transition systems of the synchronous product of these transition systems. Assume that the initial values of the registers are $r_{1}=0$ and $r_{2}=1$.
