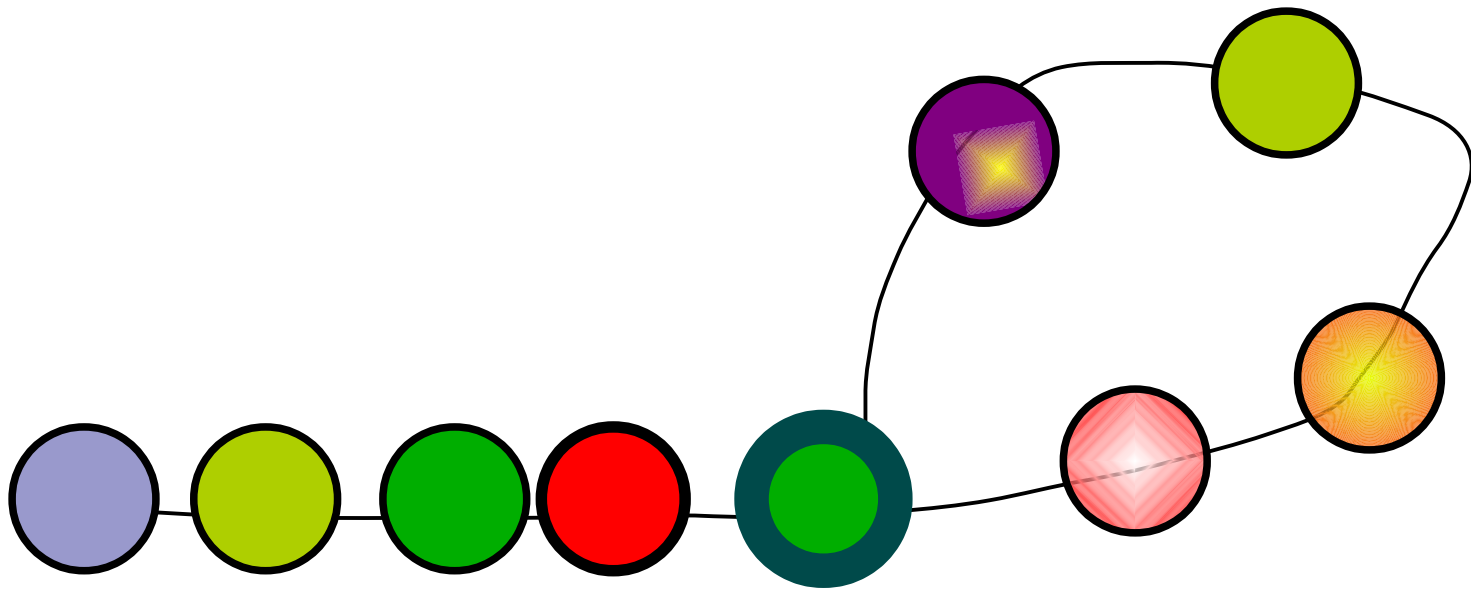
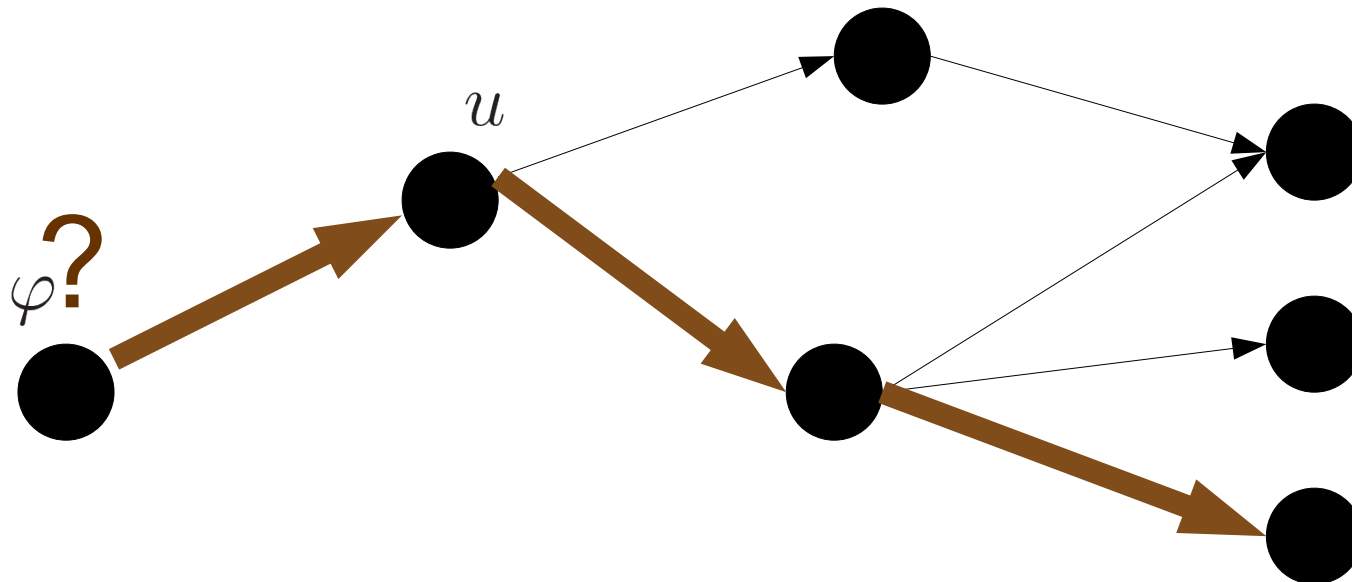


LTL-SAT



Model-checking de LTL



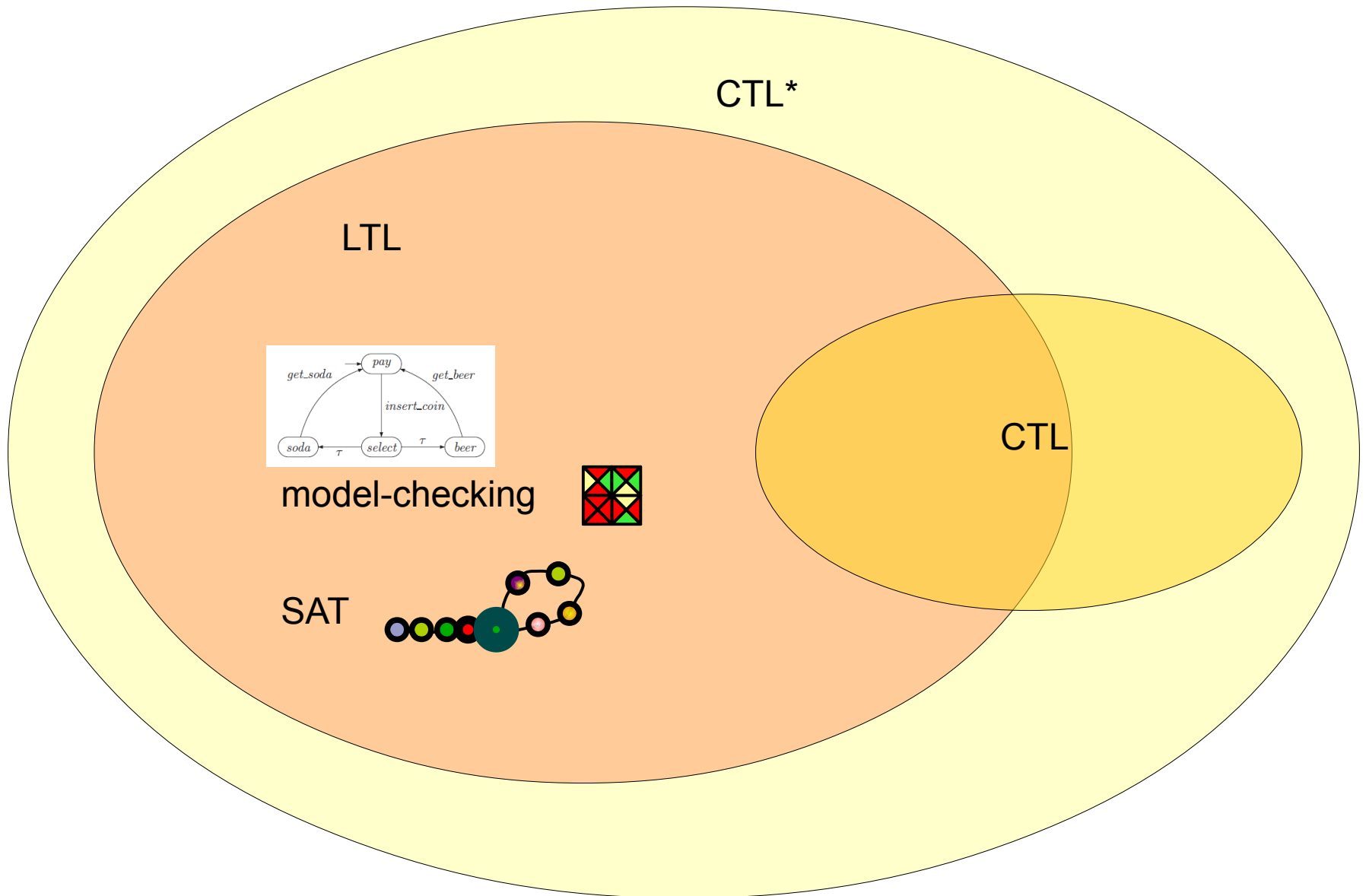
$\varphi \wedge in_w \wedge$
 $oneworld \wedge path \wedge$
 $valuations$



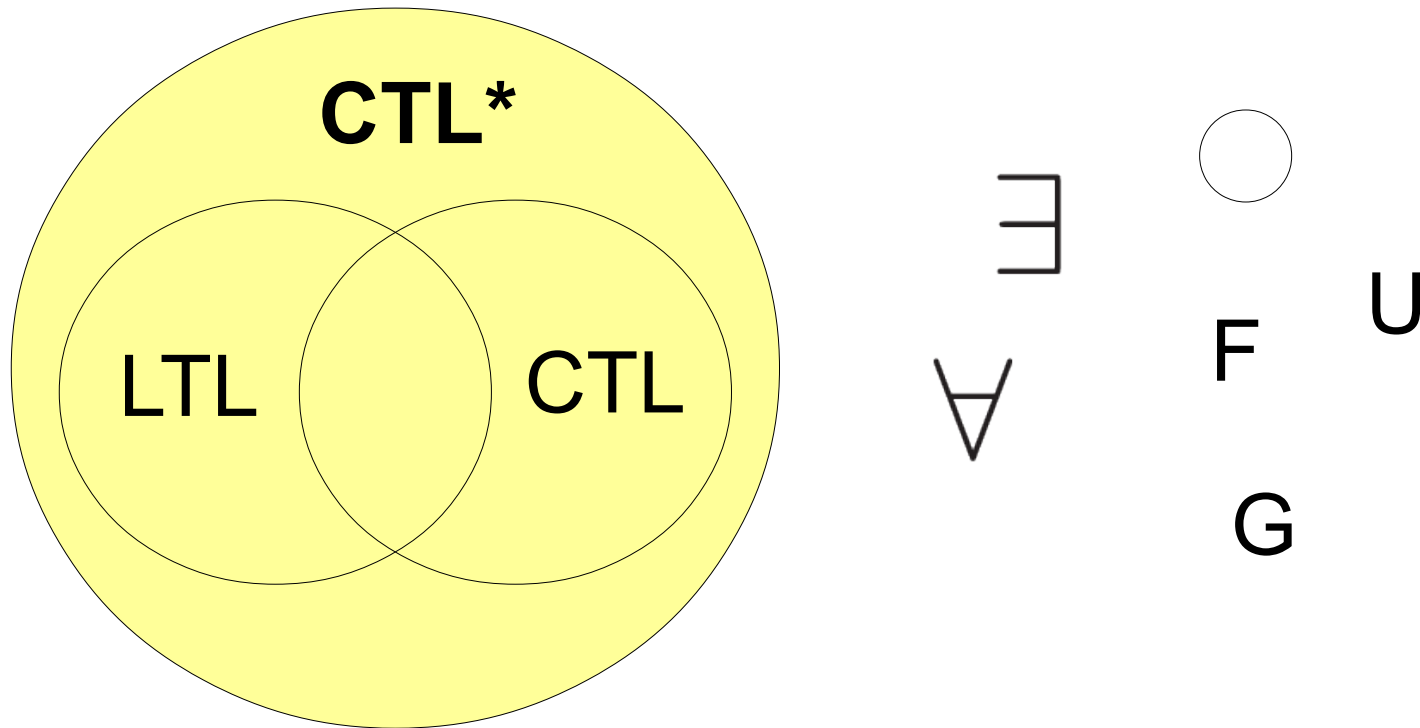
?

Retour vers CTL* : Computation Tree Logic

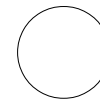
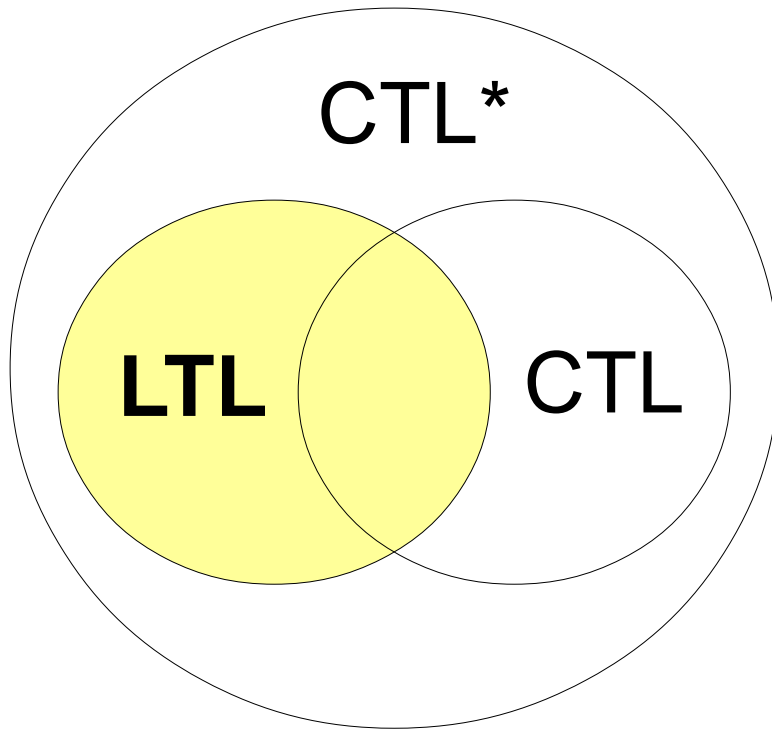
Présentation du cadre général



Computation Tree Logic*



Linear Temporal Logic

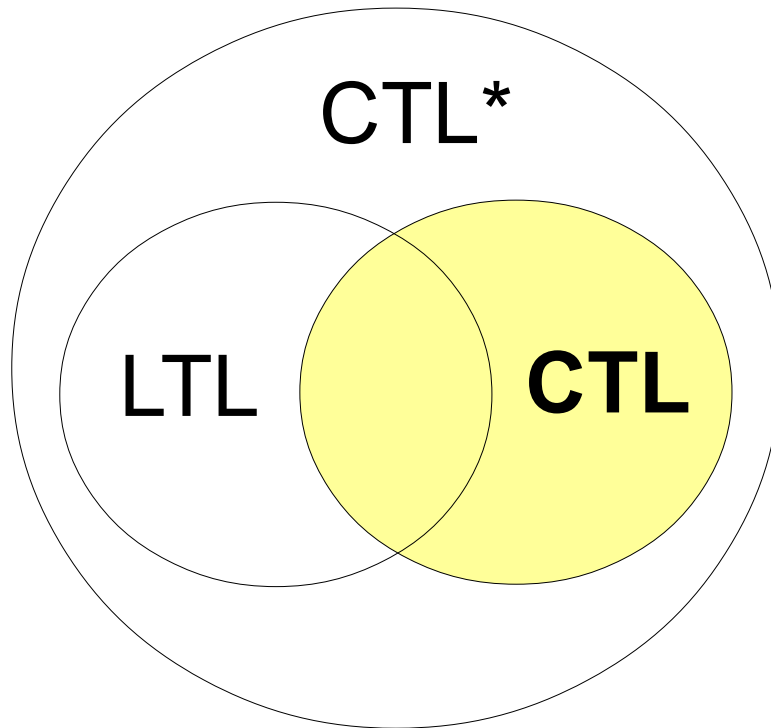


F

U

G

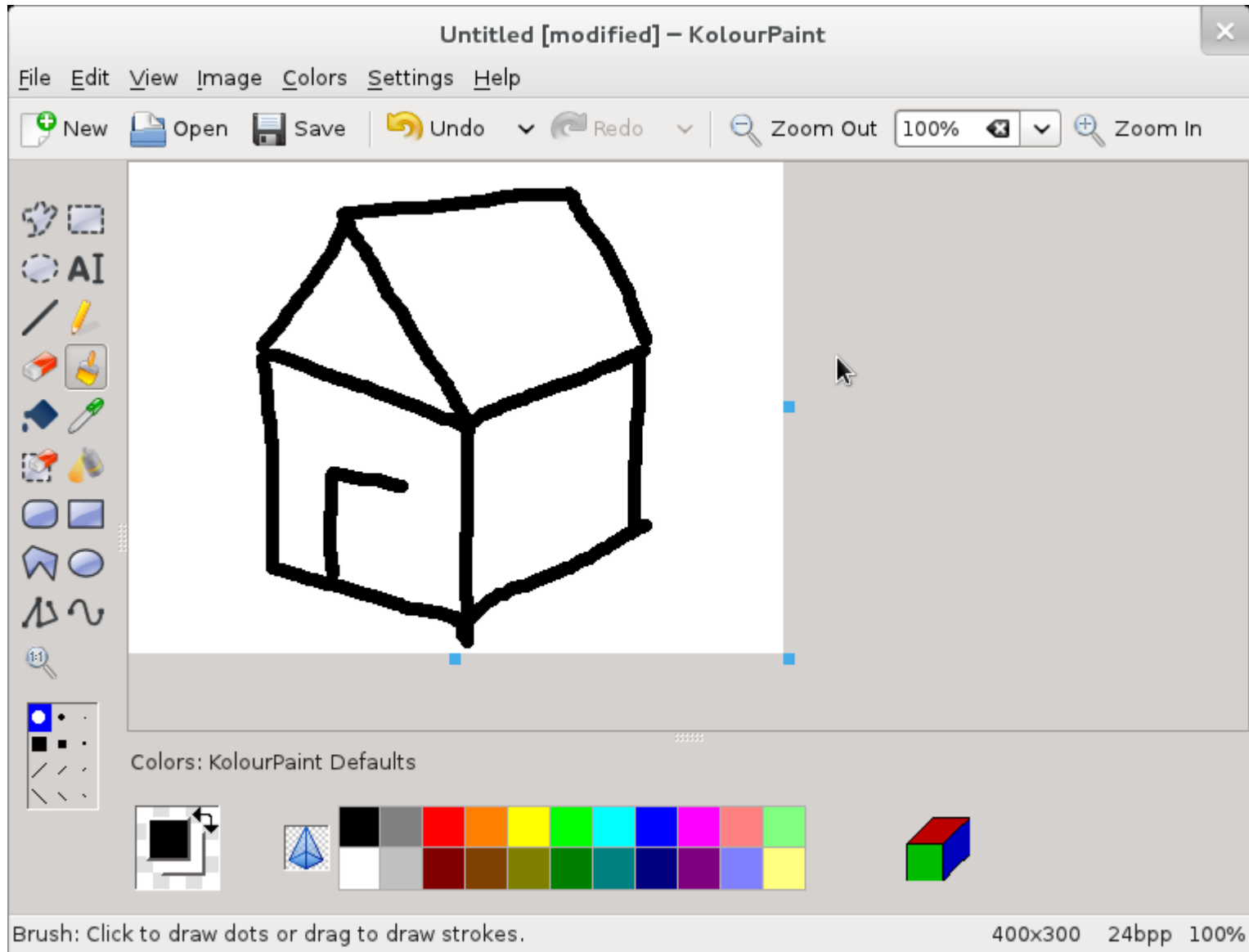
Computation Tree Logic



$\exists \bigcirc$
 $\forall \bigcirc$
 $\exists U$
 $\forall F$
 $\exists G$
 $\forall F$

Model-checking dans P !

Motivation



Propriété LTL

$\forall G(\textit{pipette} \wedge \textit{sourisappuye} \rightarrow F\neg\textit{pipette})$

$\mathcal{M}, w \models_{\forall} G(\textit{pipette} \wedge \textit{sourisappuye} \rightarrow F\neg\textit{pipette})$

Propriété LTL

$\forall G(\textit{pipette} \wedge \textit{sourisappuye} \rightarrow F\neg\textit{pipette})$

$\mathcal{M}, w \models_{\forall} G(\textit{pipette} \wedge \textit{sourisappuye} \rightarrow F\neg\textit{pipette})$

$\forall G(\textit{pipette} \wedge \textit{sourisappuye} \rightarrow \forall F\neg\textit{pipette})$

Parfois LTL ne suffit pas... on a besoin de CTL

~~$\forall G(\text{crayon} \rightarrow F \neg \text{crayon})$~~

$\exists G \text{crayon}$

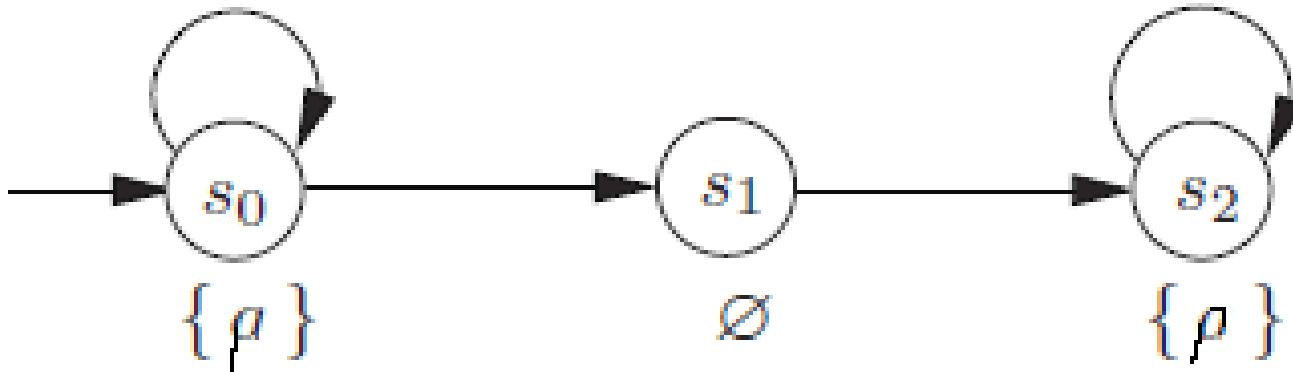
$\forall G(\text{crayon} \rightarrow \exists F \neg \text{crayon})$

$\forall G \exists F \text{crayon}$

Parfois LTL suffit mais pas CTL

$\forall FGp$

~~$\forall F\forall Gp$~~



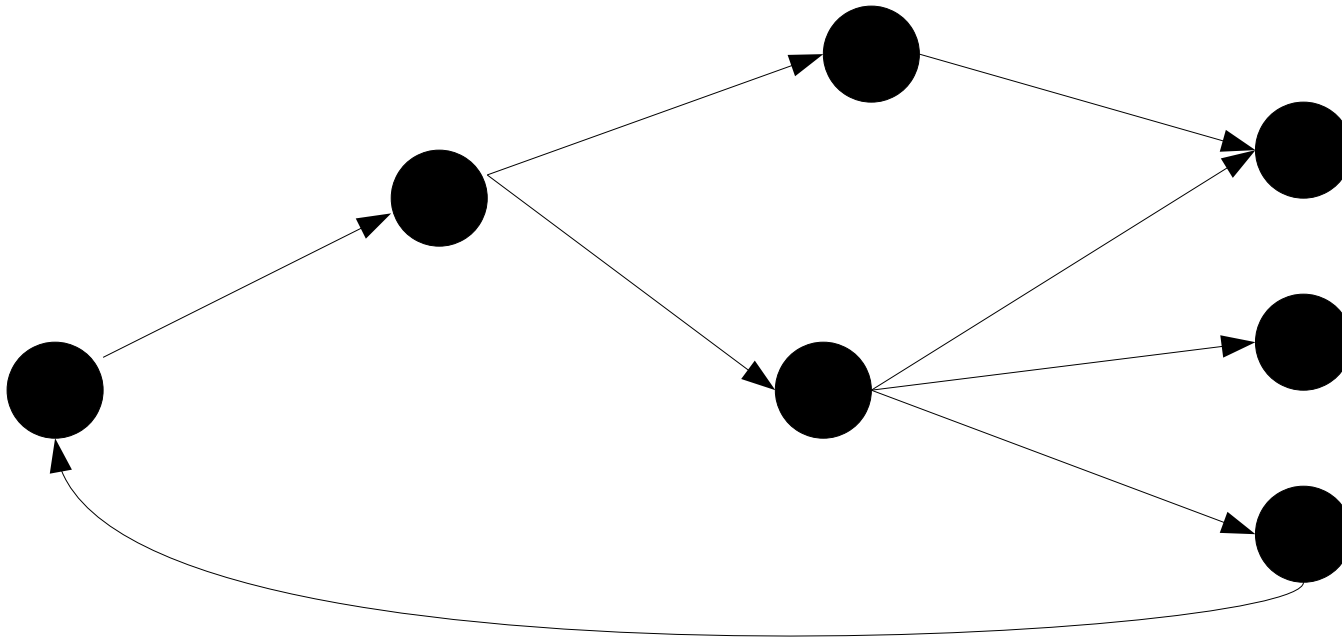
Dans de rare cas, on a besoin de CTL*

$\forall G \exists X X \text{crayon}$

$\forall G \exists X \exists X \text{crayon}$

$\forall X G \neg p \wedge \exists F G (p \vee \forall (q U p))$

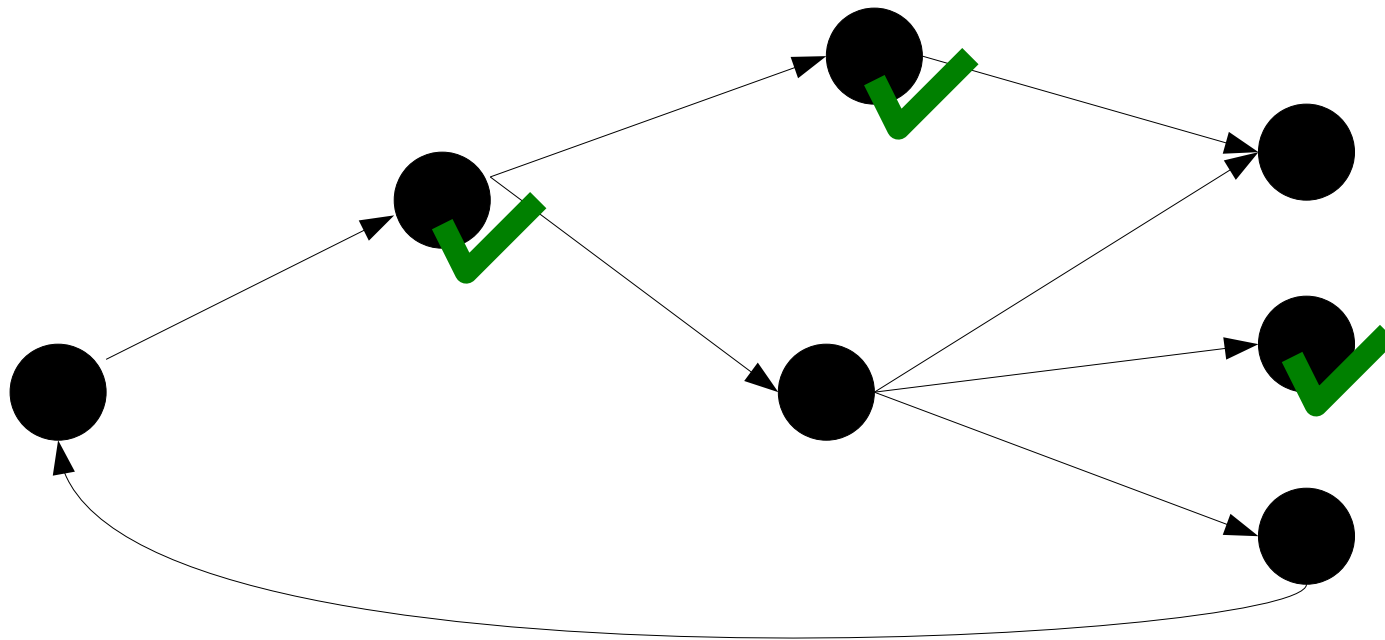
Model-checking de CTL*



$$F(p \rightarrow G\neg\exists(q \wedge G\exists FGp))$$

?

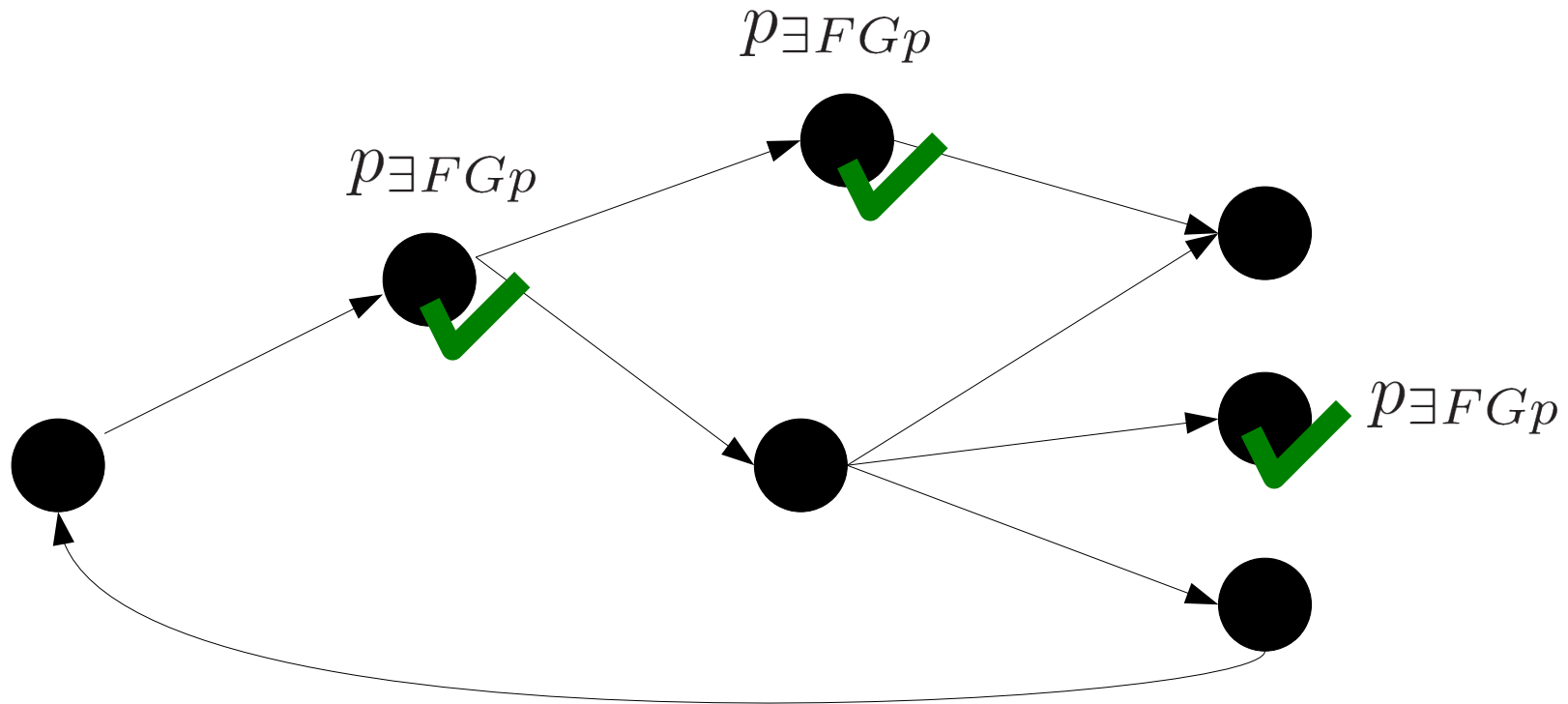
Model-checking de CTL*



$$F(p \rightarrow G\neg\exists(q \wedge G\exists FGp))$$

$$FGp \quad ?$$

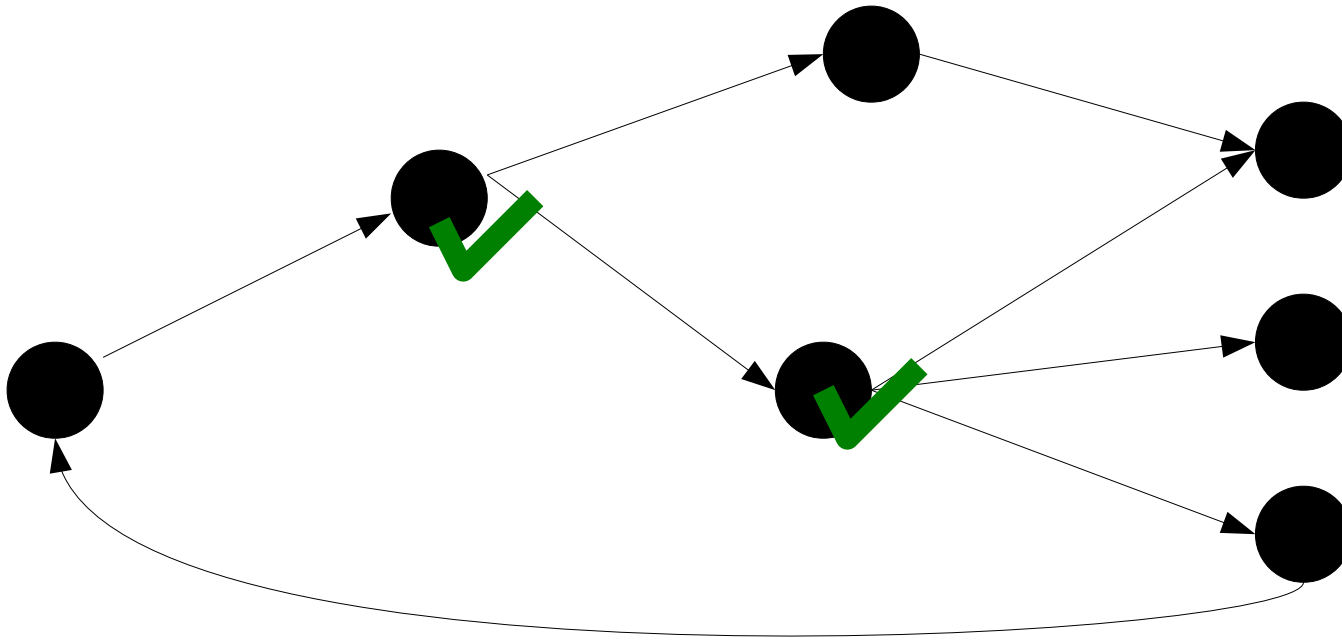
Model-checking de CTL*



$$F(p \rightarrow G \neg \exists (q \wedge G p \exists F G p))$$

?

Model-checking de CTL



$\exists F \forall (p U q)$

?

model checking CTL de $\exists\psi_1 U \psi_2$

```
 $\exists\psi_1 U \psi_2$ :  
   $S_{\psi_1} := \underline{mcCTL}(\mathcal{M}, \psi_1)$   
   $L_{totreat} := \underline{mcCTL}(\mathcal{M}, \psi_2)$   
   $result := L_{treated} := \square$   
  while  $L_{totreat} \neq \emptyset$   
     $q := L_{totreat}.defiler$   
     $result := result \cup \{q\}$   
    for  $u \rightarrow q$   
      if  $u \notin L_{treated}$  then  
         $L_{treated} := L_{treated} \cup \{u\}$   
        if  $u \in S_{\psi_1}$  then  
           $L_{totreat} := L_{totreat} \cup \{u\}$   
        endIf  
      endFor  
    endWhile  
  return  $result$ 
```

model checking CTL de $\forall \psi_1 U \psi_2$

$\forall \psi_1 U \psi_2$:

$S_{\psi_1} := \underline{mcCTL}(\mathcal{M}, \psi_1)$

for $w \in W$, $deg[w] :=$ number of successors of w

$L_{totreat} := \underline{mcCTL}(\mathcal{M}, \psi_2)$

$result := []$

while $L_{totreat} \neq \emptyset$

$q := L_{totreat}.defiler$

$result := result \cup \{q\}$

for $u \rightarrow q$

$deg[u] := deg[u] - 1$

if $deg[u] = 0$ and $u \notin result$ and $u \in S_{\psi_1}$

$L_{totreat} := L_{totreat} \cup \{u\}$

endIf

endFor

endWhile

return result