

# Compiler

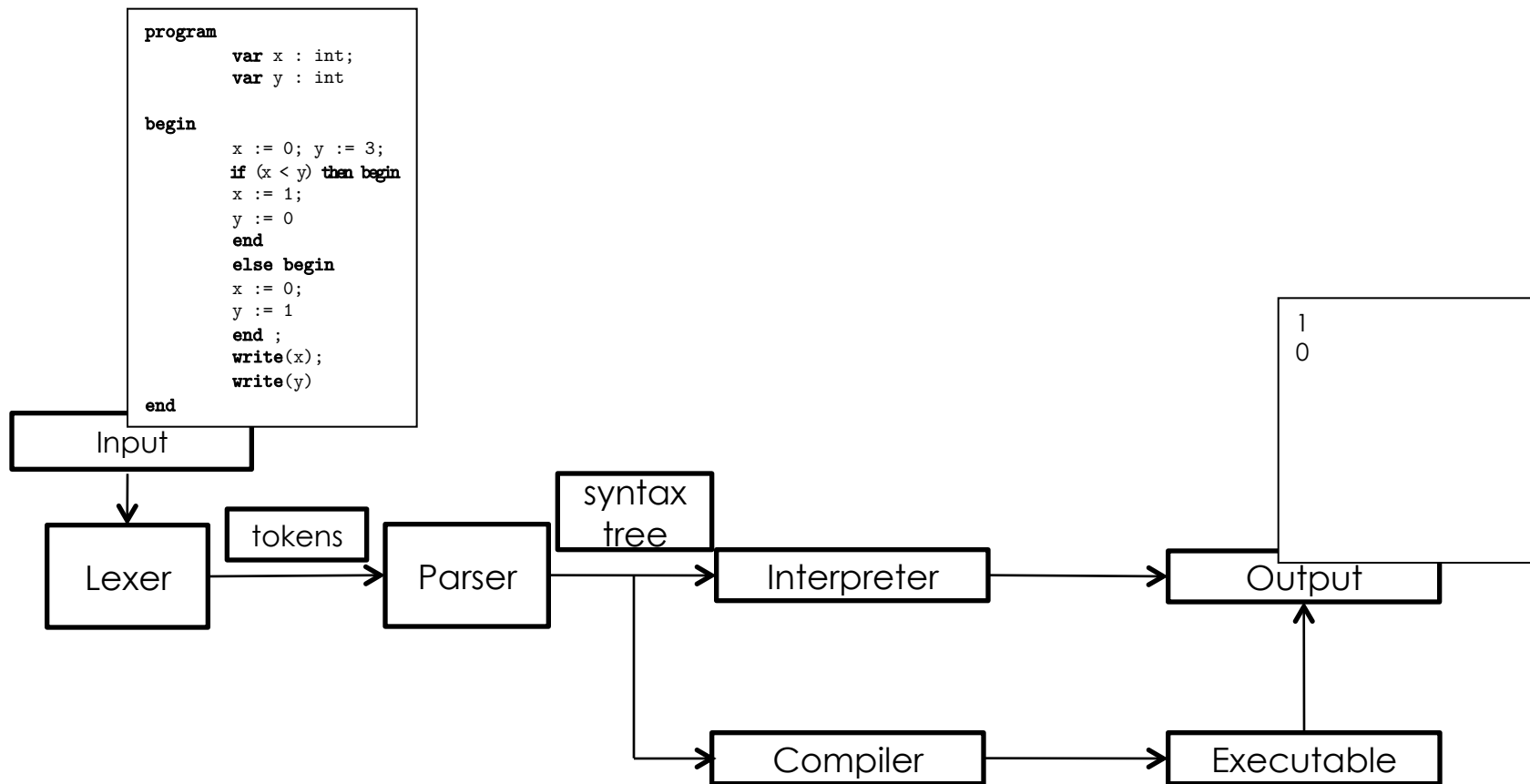
Lecture 10

Formal Languages and Compilers 2011

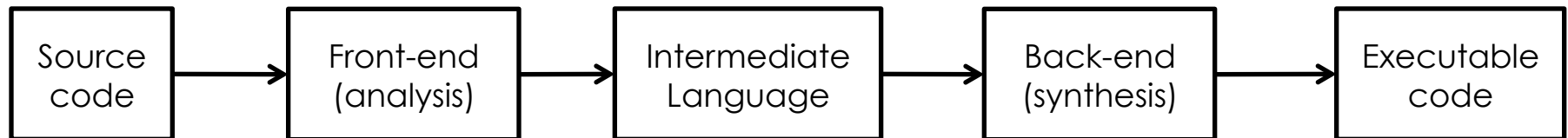
Nataliia Bielova

# Compiler: from syntax tree to target code

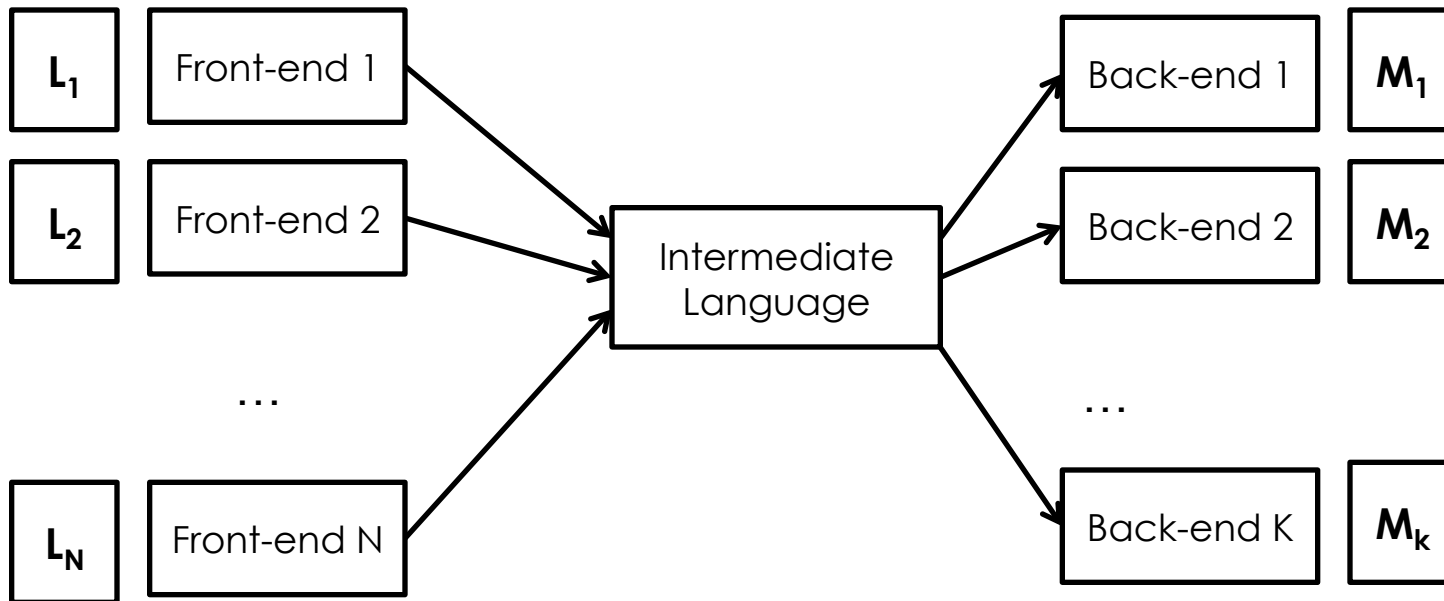
2



# Structure for a compiler

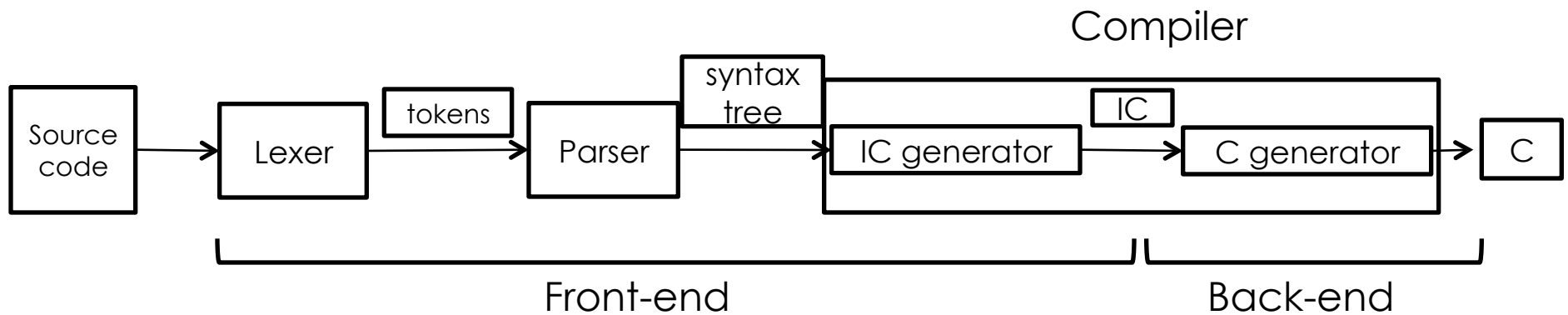


# Front-end and back-end



- Intermediate language is *platform-independent*
- Reuse the same front-end for different machines
- Reuse the same back-end for different source languages

# Compiler for crème CAraMeL



# Compiler for crème CAraMeL

Compiler:

<http://disi.unitn.it/~bielova/flc/exercises/10-Compiler.zip>

- `make`
- `./compiler < input/test_0.cre`
- [your preferred editor] `output/syntax.ast`
- [the same editor] `output/code.ic`
- [C compiler] `output/code.c`

# How the compiler is made

- `input/` file of input
- `output/` file produced by compilation
- `Documentation/index.html` documentation of the compiler
- `script.sh` script for automatic testing
- `main.ml` main program
- `lexer.mll` definition of lexical analyzer
- `parser.mly` definition of syntactic analyzer
- `syntaxtree.ml` definition of the syntax tree
- `print_syntaxtree.ml` functions for printing the syntax tree

## How the compiler is made (2)

- `exceptions.ml` exceptions used at the global level
- `declaration.ml` generalization of the table of variables (“environment”)
- `semantic.ml` type control
- `expressions.ml` translation of the evaluation of expressions
- `commands.ml` translation of the commands
- `subroutines.ml` translation of the subprograms
- `intermediate.ml` definitions of the intermediate code
- `print.ml` printing intermediate code
- `target.ml` generalization (and output) of target code



# Phases of the compiler

Syntax tree



Type control



Generation of intermediate code



Generation of target code

# Source code $\Rightarrow$ abstract syntax tree

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```
program
  var x : int

begin
  x := 3;
  write(x)
end
```

$\Rightarrow$

```
Program(
  Dec(Ide(x), Basic(Int))

Blk(
  Ass(LVar(Ide(x)), N(3))
  Write(Var(Ide(x)))
)
)
```

# Abstract syntax tree $\Rightarrow$ intermediate code

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```
Program(  
  Dec(Ide(x), Basic(Int))  
  
Blk(  
  Ass(LVar(Ide(x)), N(3))  
  Write(Var(Ide(x)))  
)  
)
```

$\Rightarrow$

CPY	Val INT: 3	NULL	offset 0
OUT	offset 0	NULL	NULL
NOP	NULL	NULL	NULL
HALT	NULL	NULL	NULL

# Intermediate language

- Is independent from the target language
- Is easy to translate to the target language (effectiveness)
- Code with 3 components:

$\underbrace{\text{NAME}}_{\text{operator}} \quad \underbrace{\text{ind}_1}_{\text{operand}_1}, \quad \underbrace{\text{ind}_2}_{\text{operand}_2}, \quad \underbrace{\text{ind}_3}_{\text{operand}_3}$

- Examples:

ADD	val <sub>1</sub>	val <sub>2</sub>	dest	
CPY	src	NULL	dest	
GOTO	label	NULL	NULL	
JNE	val <sub>1</sub>	val <sub>2</sub>	label	- Jump if Not Equal : if val <sub>1</sub> != val <sub>2</sub> goto label
CGE	val <sub>1</sub>	val <sub>2</sub>	dest	- Copy Greater or Equal: dest = (val <sub>1</sub> >= val <sub>2</sub> )
OUT	val	NULL	NULL	

# Temporal symbols

- One simple operation at a time → “cut” the expressions
- Every intermediate result should be saved in the register

```

x := (3 + 7) * 11;  ⇒  ADD    Val INT: 3  Val INT: 7    reg[2].i
                    MUL    reg[2].i  Val INT: 11   reg[1].i
                    CPY    reg[1].i  NULL         offset 0
  
```

# Labels

- Label  $\equiv$  “destination of the jump”

- Example:

	CPY	Val INT: 1	NULL	offset 0
	CPY	Val INT: 5	NULL	offset 2
	CPY	Val INT: 1	NULL	offset 1
Label2:	CGE	offset 2	offset 1	reg[1].i
	JNE	reg[1].i	Val INT: 1	Label nr. 1
	OUT	offset 1	NULL	NULL
	MUL	offset 0	offset 1	reg[2].i
	CPY	reg[2].i	NULL	offset 0
	ADD	offset 1	Val INT: 1	reg[3].i
	CPY	reg[3].i	NULL	offset 1
	NOP	NULL	NULL	NULL
	GOTO	Label nr. 2	NULL	NULL
Label1:	OUT	offset 0	NULL	NULL
	NOP	NULL	NULL	NULL
	HALT	NULL	NULL	NULL