

CVFP (Software Design and Formal Verification)

TD 1: Introduction to Software Design

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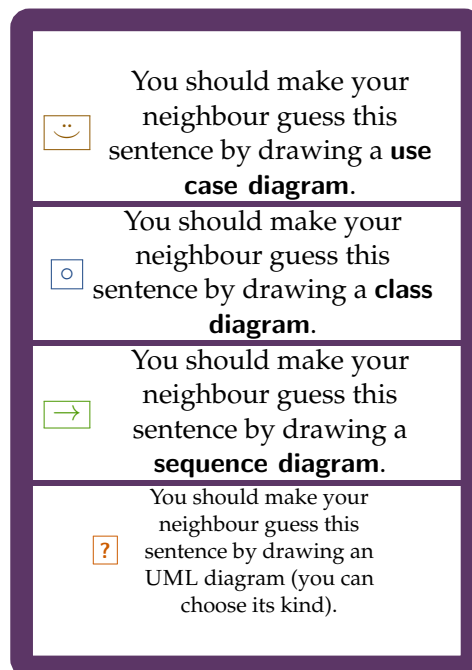
Exercise 1 As Usually In The Beginning Of The Year, We Need To Know More About Your Parents.

1.1. Design an UML object diagram representing your family tree.

1.2. Infer from it a class diagram, on which you'll annotate the various multiplicities.

Exercise 2 Bingo!

Let's apply UML while playing some PICTONARY game. You should receive soon a card with four sentences your neighbour is going to guess using UML diagrams:



Exercise 3 Morning Cross-Words

A context \mathcal{C} is a formula containing a (unique) 0-arity predicate \square named *hole*. An *instantiation* $\mathcal{C}(P)$ of such a context by a formula P is the formula $\mathcal{C}[P/\square]$, where the hole has been replaced by P .

Link every formula with a context so that every resulting formula is valid.

$\exists x, \mathcal{P}(x)$	● — ●	$\square \rightarrow \exists x, \mathcal{P}(x)$
$\forall x, (\mathcal{P}(f(x)) \wedge \neg \mathcal{P}(f(x)))$	○ ○	$\exists y, \square$
$\forall x, (\mathcal{Q}(y, y) \rightarrow \mathcal{Q}(x, x))$	○ ○	$\exists y, \square \rightarrow \forall y, \exists x, \mathcal{Q}(y, x)$
$\forall x, \mathcal{P}(x)$	○ ○	$\mathcal{P}(y) \rightarrow (\square \rightarrow \forall x, \mathcal{P}(x)) \rightarrow \mathcal{P}(f(f(y)))$
$\forall x, \mathcal{Q}(x, y)$	○ ○	$\forall x, ((\mathcal{P}(x) \rightarrow \square) \rightarrow \mathcal{P}(x)) \rightarrow \mathcal{P}(x)$
$\forall x, (\mathcal{P}(x) \rightarrow \mathcal{P}(f(x)))$	○ ○	$\neg \square \rightarrow \exists x, \neg \mathcal{P}(x)$

Exercise 4 Let's Decipher Formulae!

For each formula, give a model if it's satisfiable or (informally) explain why its negation is valid.

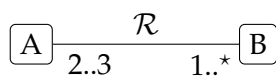
- $\forall x, \neg \mathcal{P}(x) \rightarrow \mathcal{P}(x)$
- $\forall x, \forall y, \forall z, (\mathcal{Q}(x, y) \rightarrow \mathcal{Q}(x, z) \rightarrow \mathcal{R}(y, z)) \wedge \exists x, \exists y, \neg \mathcal{R}(x, y) \wedge \forall x, \exists y, \mathcal{Q}(x, y)$
- $\forall x, \forall y, (\mathcal{Q}(x, y) \rightarrow \mathcal{Q}(y, x) \rightarrow \mathcal{R}(x, y))$
 $\wedge \forall x, \forall y, (\mathcal{Q}(x, y) \rightarrow \mathcal{Q}(y, x))$
 $\wedge \forall x, (\mathcal{Q}(x, x) \wedge \mathcal{R}(x, x))$
 $\wedge \exists x, \exists y, (\neg \mathcal{R}(x, y) \wedge \mathcal{Q}(x, y))$

Exercise 5 Never Forget That UML Diagrams Are In Fact Formulae.

5.1. Given a formula $\phi(x)$ where x is the only free-variable in $\phi(x)$, give a formula meaning that there are between 2 and 5 elements x such that $\phi(x)$.

► We can construct such a formula by splitting it into a part saying that there are at least 2 such elements and a part saying that there are at most 5.

5.2. Translate the following class diagram into a first-order formula.



5.3. Translate the class diagram designed in Exercise 1 into a first-order formula.

5.4. (*) Play the FIRST-ORDER FORMULAE Pictionary and try to succeed in at least one sentence without cheating.